

STUDY OF ALGORITHMS FOR ECHO AI VIRTUAL VOICE ASSISTANT

Prof. Londhe D.R.¹, Prasanna Raut², Samarth Sawant³, Krishna Alat⁴, Abhishek Rokade⁵

¹(Professor, SRCOE, Department of Computer Engineering Pune)

^{2,3,4,5}(Student, SRCOE, Department of Computer Engineering Pune)

Abstract: The Jarvis March algorithm, also called the Gift Wrapping algorithm, is a straightforward method used in computational geometry to determine the convex hull of a set of points. The convex hull is the smallest convex boundary that can enclose all the points, similar to the shape formed by wrapping a rubber band around them. The algorithm begins by finding the leftmost point in the set, which is guaranteed to be on the hull. From this starting point, it iteratively selects the next hull point by finding the point that is the most counterclockwise relative to the current point. This process continues, effectively "wrapping" around the point set, until the algorithm returns to the starting point, completing the convex hull. The Jarvis March algorithm, also called the Gift Wrapping algorithm, is a straightforward method used in computational geometry to determine the convex hull of a set of points. The convex hull is the smallest convex boundary that can enclose all the points, similar to the shape formed by wrapping a rubber band around them. The algorithm begins by finding the leftmost point in the set, which is guaranteed to be on the hull. From this starting point, it iteratively selects the next hull point by finding the point that is the most counterclockwise relative to the current point. This process continues, effectively "wrapping" around the point set, until the algorithm returns to the starting point, completing the convex hull.

Key Word: Voice Assistant, Natural Language Processing, Desktop Assistant, Speech Recognition

I. INTRODUCTION

The convex hull of a set of points (S) is the smallest convex polygon that contains all of the points in S [1]. Computing convex hulls is a fundamental issue with a wide range of applications in computer sciences and mathematics. Some of the applications in which the convex hull is used include image processing, pattern recognition, medical simulations, geographical information systems (GIS), pathfinding, and computer visualization [2]. It is also important for video games, simulations of many bodies physically interacting, engineering design, and recently, in autonomous driving, where computations must be strictly fast and in real time [3]. Many classic algorithms have been introduced for computing convex hulls, including the Graham scan (1972) [4], Jarvis's march (1973) [5], the divide-and-conquer algorithm (1977) [6], Andrew's monotone chain (Andrew 1979) [7], the incremental approach (1984) [8], and QuickHull (1996) [9]. However, one aspect can further improve the performance of these algorithms: adding a preprocessing stage to reduce computation time and memory space. One of the most widely used algorithms for finding convex hulls is the Graham scan algorithm, which calculates the convex hull of a set of points in $O(n \log n)$.

Author suggested solution presented in this research is to enhance the Graham scan algorithm [5] by adding filtering techniques for fast convex hull computing. The important stage that is added to Graham's algorithm is discarding points from the approximate convex hull and inserting them into priority queues. The proposed algorithm was designed for serial computing, so it does not need any graphics processing unit (GPU) resources. However, the challenge is to find an data.

II. LITERATURE REVIEW

Jayaram, M., Fleyeh, H. (2016) "Convex Hulls in Image Processing A convex hull algorithm is a fundamental computational geometry technique used to find the smallest convex boundary that encompasses a set of points on a plane, similar to stretching a rubber band around them. The convex hull, a polygon formed by the outermost points, can be computed using various algorithms, each suited to different scenarios based on efficiency and dataset size. For instance, the Jarvis March algorithm (Gift Wrapping) starts from the leftmost point and iteratively selects the point with the most counterclockwise angle relative to the current point until it returns to the starting position, with a

complexity of $O(n \cdot h)O(n \cdot h)$, where h is the number of points on the hull. Graham's Scan algorithm, more efficient for larger sets, first sorts points by polar angle from a pivot, then builds the hull by checking turns between points, achieving a complexity of $O(n \log n)O(n \log n)$. Another approach, QuickHull, resembles QuickSort by dividing points with a baseline and recursively identifying extreme points, performing efficiently in practice at $O(n \log n)O(n \log n)O(n \log n)$ on average, though it can reach $O(n^2)O(n^2)O(n^2)$ in the worst case.

Mei, G. (2016). "CudaChain: an alternative algorithm for finding 2D convex hulls on the GPU." SpringerPlus, 5(1). Jarvis March algorithm (Gift Wrapping) starts from the leftmost point and iteratively selects the point with the most counterclockwise angle relative to the current point until it returns to the starting position, with a complexity of $O(n \cdot h)O(n \cdot h)O(n \cdot h)$, where h is the number of points on the hull. Graham's Scan algorithm, more efficient for larger sets, first sorts points by polar angle from a pivot, then builds the hull by checking turns between points, achieving a complexity of $O(n \log n)O(n \log n)O(n \log n)$. Another approach, QuickHull, resembles QuickSort by dividing points with a baseline and recursively identifying extreme points, performing efficiently in practice at $O(n \log n)O(n \log n)O(n \log n)$ on average, though it can reach $O(n^2)O(n^2)O(n^2)$ in the worst case. The Monotone Chain algorithm, a variant of Graham's Scan, sorts points lexicographically and constructs the hull by iterating twice through the sorted points to form upper and lower hulls. Each method has unique strengths, such as QuickHull's speed in practice or Graham's Scan's reliability on large datasets.

Graham, RL. "An efficient algorithm for determining the convex hull of a finite planar set." Inform Process Lett 1972; 1(4). The Jarvis March algorithm, also known as the Gift Wrapping algorithm, is a straightforward method in computational geometry used to find the convex hull of a set of points in a 2D plane. The convex hull is the smallest convex boundary that can enclose all points in a dataset, similar to wrapping a rubber band around the points. Jarvis March is intuitive and works well for small datasets, though it can be less efficient for larger datasets compared to other convex hull algorithms. Jarvis March is commonly used in computer graphics, geographic information systems, and robotics. It's useful in scenarios where the convex hull is needed for small datasets or where computational simplicity is preferred over speed.

Barber CB, Dobkin DP, Huhdanpaa H. "The quickhull algorithm for convex hulls." ACM Transactions on Mathematical Software (TOMS), Volume 22 Issue 4, Dec. 1996, Pages 469-483. The Convex Hull algorithm is a foundational technique in computational geometry, designed to compute the smallest convex boundary that fully encloses a set of points in a 2D plane, resembling a rubber band stretched around the outermost points. This convex boundary, or hull, can be found using various algorithms, each with specific strengths depending on the dataset. For example, Jarvis March (Gift Wrapping) is an intuitive method that iteratively selects the most counterclockwise point relative to the current hull point, giving it a complexity of $O(n \cdot h)O(n \cdot h)O(n \cdot h)$, where h is the number of points on the hull. This makes it suitable for smaller datasets but less efficient for large ones. In contrast, Graham's Scan and Monotone Chain both start by sorting points based on angles or coordinates and then iteratively build the hull by ensuring only counterclockwise turns, achieving an overall complexity of $O(n \log n)O(n \log n)O(n \log n)$ due to the sorting step, which makes them more efficient for larger datasets. QuickHull, inspired by QuickSort, divides points relative to a baseline and recursively identifies extreme points, performing efficiently in practice at $O(n \log n)O(n \log n)O(n \log n)$ on average. Convex hull algorithms are crucial in fields like computer graphics, geographic information systems, robotics, and image processing, where defining the boundaries of a point set helps in tasks like collision detection, shape analysis, and pathfinding. Each algorithm's unique approach makes it adaptable to different datasets and requirements, ensuring the fastest and most accurate boundary calculation based on the input characteristics.

Qin, J., Mei, G., Cuomo, S., Guo, S. and Li, Y. (2019). "CudaCHPre2D: A straightforward preprocessing approach for accelerating 2D convex hull computations on the GPU." Concurrency and Computation: Practice and Experience, p.e5229. Convex hull algorithms are crucial in fields like computer graphics, geographic information systems, robotics, and image processing, where defining the boundaries of a point set helps in tasks like collision detection, shape analysis, and pathfinding. Each algorithm's unique approach makes it adaptable to different datasets and requirements, ensuring the fastest and most accurate boundary calculation based on the input characteristics. This process continues, effectively "wrapping" around the point set, until it returns to the starting point, completing the convex hull. Jarvis March has a time complexity of $O(n \cdot h)O(n \cdot h)O(n \cdot h)$, where n is the total number of points and h is the number of points on the hull, making it simple but inefficient for large datasets. The algorithm is often used in scenarios with smaller datasets or when only a few points are expected on the hull, such as in computer graphics, geographic mapping, and boundary detection.

III . Jarvis and Convex Hull Algorithm

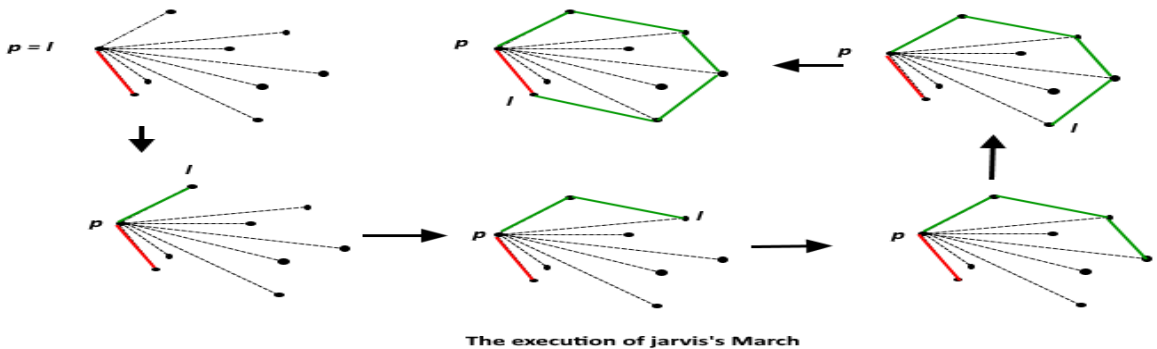
1. Jarvis March Algorithm:

The Jarvis March algorithm, also known as the Gift Wrapping algorithm, is a method for computing the convex hull of a set of points in 2D space. The convex hull of a set of points is the smallest convex polygon that can enclose all the points. Jarvis March is particularly effective when the number of points on the convex hull is small, as it runs in $O(nh)O(nh)O(nh)$ time complexity, where n is the total number of points, and h is the number of points on the hull. Consider the general case when the input to the algorithm is a finite unordered set of points on a Cartesian plane. An important special case, in which the points are given in the order of traversal of a simple polygon's boundary, is described later in a separate subsection. If not all points are on the same line, then their convex hull is a convex polygon whose vertices are some of the points in the input set. Its most common representation is the list of its vertices ordinary.

Steps of Jarvis March Algorithm:

1. Identify the leftmost point (or, in some variations, the point with the smallest y -coordinate). This point is guaranteed to be on the convex hull, as it cannot be enclosed by other points.
2. Iteratively add points to the hull by selecting the point that makes the smallest counterclockwise angle with the line from the last hull point. This ensures we "wrap around" the outer boundary in a counterclockwise direction.
3. Terminate the process when we return to the starting point, completing the loop around the convex hull.

Architecture of Jarvis March Algorithm :



The execution of jarvis's March

Fig 1 . Jarvis March Algorithm

Advantages of Jarvis March Algorithm :

1. **Personalization** : ECHO learns from user interactions, preferences, and behaviors over time in This personalization allows it to offer tailored responses, reminders, and suggestions, making the experience more relevant and engaging for each individual
2. **Multitasking Capabilities**: ECHO can handle multiple requests simultaneously, such as setting reminders, playing music, and providing weather updates. This multitasking ability enhances efficiency, allowing users to accomplish more in less time.
3. **Smart Home Integration** : ECHO is compatible with a wide range of smart home devices, enabling users to control lights, thermostats, and security systems with voice commands. This integration creates a cohesive smart home environment, enhancing convenience and energy management.
4. **Accessibility** : ECHO provides hands-free assistance, making it an invaluable tool for individuals with disabilities or mobility challenges. Voice commands allow these users to access information and control devices without needing to use traditional interfaces.
5. **Structured and Efficient Decision Making**: The decision tree approach allows for systematic and quick checks against ABO and Rh compatibility rules, making it computationally efficient for binary classification decisions.

Disadvantages of Jarvis March Algorithm :

1. Privacy Concerns : ECHO continuously listens for voice commands, which raises concerns about user privacy. Despite security measures, there is a risk of sensitive information being inadvertently recorded or misused, leading to hesitation among some users.

2. Dependence on Internet Connectivity : ECHO requires a stable internet connection to function effectively. In areas with poor connectivity, the assistant's capabilities can be significantly limited, impacting user experience and accessibility.

3. Limited Understanding of Context : While ECHO has contextual awareness, it can still struggle with complex or ambiguous queries. It may misinterpret user intent, leading to irrelevant or inaccurate responses, which can be frustrating for users.

4. Voice Recognition Limitations : ECHO's performance can be affected by background noise, multiple speakers, or heavy accents. This can lead to difficulties in accurately recognizing commands, reducing the efficiency of interactions.

Applications of Jarvis March Algorithm :

1. In GIS, convex hulls are used to determine the outer boundary of geographical features, like forested areas, lakes, or urban regions, based on GPS coordinates. The Jarvis March algorithm can help to efficiently outline the perimeter of these regions when the hull size is relatively small.

2. Convex hulls are used in computer graphics for object boundary detection, collision detection, and rendering of complex shapes. Jarvis March can be useful when working with a small number of edge points, such as creating bounding boxes for sprites or objects in 2D games.

2. Convex Hull Algorithm :

A convex hull of a set of points in a 2D plane is the smallest convex polygon that encloses all the points. Convex hull algorithms find this polygon by identifying a subset of the input points that form the "outer shell" or boundary. Given a set of points in a plane, the convex hull is the minimal convex boundary that completely contains all the points. Imagine stretching a rubber band around a set of nails hammered into a board, where the rubber band snaps to form a boundary around the outermost nails — this boundary is the convex hull. The Convex Hull problem is a fundamental problem in computational geometry. Given a set of points in a plane, the task is to find the smallest convex polygon that can enclose all of the points, effectively forming a "convex boundary" around them. There are several algorithms for solving this, including Graham's Scan, Jarvis March, and Quick hull, each with different characteristics and complexities. (NLP techniques, allowing it to understand and respond to user queries in a conversational manner. This enables users to interact with the assistant naturally, as they would with another person, reducing the learning curve for new users. The next stage of the algorithm is a complete pass on the input S . For each point $p_i \in S$, the algorithm checks if it is inside the polygon or not. If the current point lies inside the polygon, then this point is filtered by the algorithm. Importantly, in many practical cases where the distribution function of the input set S is similar to the normal or uniform distribution, most of the points are inside of the filtering polygon and will be discarded in linear time. If the point lies outside P , then the algorithm will insert the point into priority queues, which are implemented using min/max heaps. The queue Q holds the points of the four regions formed by P and the corner points.

Steps of Convex Hull Algorithm:

1. Identify the leftmost point : Find the point with the smallest x coordinate (and lowest y coordinate in case of tie). This point is guaranteed to be part of the convex hull and serves as the starting point.
2. Initialize the Hull : Set the starting point as the current hull point
3. Find the next hull point : If a point forms a counterclockwise angle or is farthest in a colinear alignment, update the next hull point.
4. Add to Hull and Repeat : Repeat this process until returning to the starting point, completing the convex hull.
5. Repeat the process to follow the sequence

Architecture of Convex Hull Algorithm :

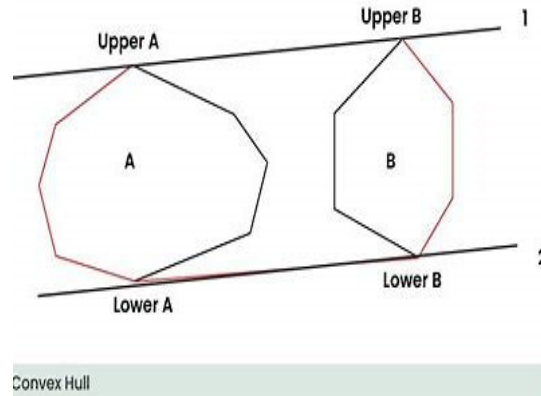


Fig 2 . Convex Hull Algorithm

Advantages of Convex Hull Algorithm :

- 1.Contextual Awareness :** ECHO utilizes contextual understanding to provide more relevant answers based on previous interactions and the current environment. This capability enhances the quality of responses, making them more useful and timely.
- 2. Continuous Learning :** Through machine learning, ECHO continuously improves its understanding and processing of language. As it interacts with users, it becomes more adept at handling complex queries and adapting to changing user needs.
- 3. Wide Range of Applications :** ECHO can assist with a variety of tasks, from managing calendars

Disadvantages of Convex Hull Algorithm :

- 1. Voice Recognition Limitations :** ECHO's performance can be affected by background noise, multiple speakers, or heavy accents. This can lead to difficulties in accurately recognizing commands, reducing the efficiency of interactions.
- 2. Potential for Over-Reliance :** Users may become overly dependent on ECHO for tasks and information, which can diminish their ability to perform these tasks independently. This reliance may hinder the development of problem-solving skills.
- 3. Security Vulnerabilities :** Like any connected device, ECHO is susceptible to hacking and cyber threats. If not properly secured, personal data and home networks could be compromised.
- 4. Learning Curve for Some Users :** Although ECHO is designed to be user-friendly, individuals who are not tech-savvy may still face a learning curve. Understanding how to utilize all features effectively can be challenging for some users.

Applications of Convex Hull Algorithm :

1. Media Control : Play music, videos, or podcasts, and adjust audio settings across devices
2. Reading Assistance : Read books aloud, provide summaries, and quiz users on content
3. Convert spoken words into written text for people with mobility challenges

Conclusion

The Jarvis March algorithm provides an effective solution to the convex hull problem in computational geometry, particularly for datasets where the convex hull includes only a fraction of the total points. It remains relevant due to its simple, angle-based approach, which is both intuitive and beneficial in cases with sparse point distributions. Despite being outperformed by more complex algorithms for large datasets or densely populated hulls, it remains a viable choice for problems with specific point distributions, making it an algorithm worth understanding and implementing when such datasets arise. The field of Natural Language Processing (NLP) encompasses a wide range of algorithms that enable machines to process, analyze, understand, and generate human language. The conclusion for NLP algorithms is multifaceted, reflecting advancements in machine learning, deep learning, and linguistic theory, as well as practical challenges and ongoing research areas. Here's a detailed breakdown of the conclusion for NLP algorithms

Reference

- [1] Jayaram, M., Fleyeh, H. (2016) "Convex Hulls in Image Processing: A Scoping Review." *American Journal of Intelligent Systems* 6 (2): 48-58.
- [2] Mei, G. (2016). "CudaChain: an alternative algorithm for finding 2D convex hulls on the GPU." *SpringerPlus*, 5(1).
- [3] SOUTO, N. (2019). "Video Game Physics Tutorial - Part II: Collision Detection for Solid Objects." [online] DEVELOPERS.
- [4] Graham, RL. "An efficient algorithm for determining the convex hull of a finite planar set." *Inform Process Lett* 1972; 1(4).
- [5] Jarvis, RA. "On the identification of the convex hull of a finite set of points in the plane." *Inform Process Lett* 1973; 2(1)
- [6] Franco P. Preparata. S.J. Hong. "Convex Hulls of Finite Sets of Points in Two and Three Dimensions," *Commun. ACM*, vol. 20, no. 2, pp. 87-93, 1977.
- [7] A.M. Andrew. "Another efficient algorithm for convex hulls in two dimensions." *Information Processing Letters*, Volume 9, Issue 5, 16 December 1979, Pages 216-219.
- [8] Kallay, M. "The complexity of incremental convex hull algorithms in R^d ." *Information Processing Letters*, Volume 19, Issue 4, 12 November 1984, Page 197. [9] Barber CB, Dobkin DP, Huhdanpaa H. "The quickhull algorithm for convex hulls." *ACM Transactions on Mathematical Software (TOMS)*, Volume 22 Issue 4, Dec. 1996, Pages 469-483.
- [9] Stein, A., Geva, E. and El-Sana, J. (2012). "CudaHull: Fast parallel 3D convex hull on the GPU." *Computers & Graphics*, 36(4), pp.265-271
- [10] Qin, J., Mei, G., Cuomo, S., Guo, S. and Li, Y. (2019). "CudaCHPre2D: A straightforward preprocessing approach for accelerating 2D convex hull computations on the GPU." *Concurrency and Computation: Practice and to imExperience*, p.e5229.
- [11] Tang, M., Zhao, J., Tong, R. and Manocha, D. (2012). "GPU accelerated convex hull computation." *Computers & Graphics*, 36(5), pp.498-506. [13] Gao, M., Cao, T., Nanjappa, A., Tan, T. and Huang, Z. (2013). "gHull.: A GPU algorithm for 3D convex hull." *ACM Transactions on Mathematical Software*, 40(1), pp.1-19.
- [12] Neelam Labhade-Kumar "Voice operated assistant system for blind people Using Machine Learning", *Journal of Harbin Engineering University*, Scopus, Issue 3 volume 45, PP 191-197, march 2024
- [13] Neelam Labhade-Kumar "To Study Different Types of Supervised Learning Algorithm" May 2023, *International Journal of Advanced Research in Science, Communication and Technology (IJARSCT)*, Volume 3, Issue 8, May 2023, PP-25-32, ISSN-2581-9429 DOI: 10.48175/IJARSCT-10256
- [14] Neelam L-Kumar "Developing interpretable models and techniques for explainable AI in decision-making", *The Scientific Temper* (2023) UGC Care-II Vol. 14 (4): 1324-1331, E-ISSN: 2231-6396, ISSN: 0976-8653, Published : December 2023